

A Simple *Non-Interactive* Pairwise Key Establishment Scheme in Sensor Networks

Chia-Mu Yu^{1,2}, Chun-Shien Lu¹, and Sy-Yen Kuo²

¹*Academia Sinica*

²*National Taiwan University*

Overview

- Sensor network
 - A lot of resource-constrained sensor nodes
 - Adversarial environment
- Goal of this paper
 - Establishing key between any pair of nodes
- Node compromise

Key Establishment in Sensor Network

- **Numerous** proposals
 - First work is [CCS'02]
 - Then, [S&P'03], [CCS'03], [INFOCOM'04], [INFOCOM'05], [INFOCOM'06], [MobiHoc'07], etc
 - 872 results by [Google Scholar](#)
- Why one more scheme?
 - **Interactive** becomes **non-interactive**

Why *interactive* matters?

- Energy

- Communication dominates the energy consumption in sensor networks[MobiCom'01]
- reducing #comm. prolongs network lifetime

- Security

- DoS attack which sends bogus key-sharing message to distant node

What We Want?

- **Desired properties**

1. Resilience to the Adversary's Intervention

- Eavesdropping, node compromise, and PHY/MAC attacks

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1. Resilience to the Adversary's Intervention
 - Eavesdropping, node compromise, and PHY/MAC attacks
2. Directed and Guaranteed Key Establishment
 - Node compromise

What We Want?

- **Desired properties**

1. Resilience to the Adversary's Intervention
 - Eavesdropping, node compromise, and PHY/MAC attacks
2. Directed and Guaranteed Key Establishment
 - Mobile sink, node compromise
3. Resilience to Network Configurations
 - Mobility, heterogeneity, hardware, etc

What We Want?

- **Desired properties**

4. Efficiency

- Computation, communication, and storage

What We Want?

- **Desired properties**

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5. Resilience to Dynamic Node Deployment

- Re-deploy some sensor nodes

What We Want?

- **Desired properties**

4. Efficiency

- Computation, communication, and storage

5. Resilience to Dynamic Node Deployment

- Re-deploy some sensor nodes

6. Transparency

- Power-saving protocols

Related Work

- [EuroCrypt'84] $A = (D \cdot G)^T$ and $K = A \cdot G$

$$\begin{array}{c}
 A \\
 \begin{array}{c} A_{1,-} \\ A_{3,-} \end{array}
 \end{array}
 \begin{array}{|c|c|c|c|}
 \hline
 21 & 67 & 48 & 88 \\
 \hline
 24 & 66 & 52 & 90 \\
 \hline
 22 & 19 & 29 & 41 \\
 \hline
 26 & 67 & 55 & 93 \\
 \hline
 \end{array}
 \cdot
 \begin{array}{c}
 G \\
 \begin{array}{c} G_{-1} \\ G_{-3} \end{array}
 \end{array}
 \begin{array}{|c|c|c|c|}
 \hline
 5 & 1 & 15 & 1 \\
 \hline
 11 & 5 & 0 & 4 \\
 \hline
 7 & 9 & 1 & 0 \\
 \hline
 9 & 14 & 6 & 15 \\
 \hline
 \end{array}
 =
 \begin{array}{c}
 K \\
 \begin{array}{c} K_{3,1} \\ K_{1,3} \end{array}
 \end{array}
 \begin{array}{|c|c|}
 \hline
 891 \\
 \hline
 891 \\
 \hline
 \end{array}$$

- [MobiHoc'07]
 - RPB scheme
 - Random perturbation

Idea

$$\begin{array}{c}
 W \\
 W_{1,-} \\
 W_{3,-}
 \end{array}
 \begin{bmatrix}
 20 & 68 & 48 & 88 \\
 24 & 66 & 52 & 90 \\
 21 & 18 & 29 & 42 \\
 26 & 67 & 55 & 93
 \end{bmatrix}
 \cdot
 \begin{array}{c}
 G \\
 G_{-,1} \\
 G_{-,3}
 \end{array}
 \begin{bmatrix}
 5 & 1 & 15 & 1 \\
 11 & 5 & 0 & 4 \\
 7 & 9 & 1 & 0 \\
 9 & 14 & 6 & 15
 \end{bmatrix}
 =
 \begin{array}{c}
 K \\
 K'_{3,1} \\
 K'_{1,3}
 \end{array}
 \begin{bmatrix}
 884 \\
 876
 \end{bmatrix}$$

$$[20 \ 68 \ 48 \ 88] = W_{1,-} = A_{1,-} + \phi_1 = [21 \ 67 \ 48 \ 88] + [-1 \ 1 \ 0 \ 0]$$

$$[21 \ 18 \ 29 \ 42] = W_{3,-} = A_{3,-} + \phi_3 = [22 \ 19 \ 29 \ 41] + [-1 \ -1 \ 0 \ 1]$$

ϕ_1 and ϕ_3 are the random noise for $W_{1,-}$ and $W_{3,-}$, respectively

$$K'_{1,3} \neq K'_{3,1} \quad \text{but} \quad
 \begin{array}{l}
 X_{1,3} = (11011)_2 = f_{10,5}(876) \\
 X_{3,1} = (11011)_2 = f_{10,5}(884)
 \end{array}$$

If represented by 10 bits, their first 5 bits will be the same ¹²

Idea

- Perturbed result should be under control

$$\begin{array}{c}
 W \\
 \begin{array}{c}
 W_{1,-} \\
 W_{3,-}
 \end{array}
 \end{array}
 \begin{bmatrix}
 20 & 68 & 48 & 88 \\
 24 & 66 & 52 & 90 \\
 21 & 18 & 29 & 42 \\
 26 & 67 & 55 & 93
 \end{bmatrix}
 \cdot
 \begin{array}{c}
 G \\
 \begin{array}{c}
 G_{-1} \\
 G_{-3}
 \end{array}
 \end{array}
 \begin{bmatrix}
 5 & 1 & 15 & 1 \\
 11 & 5 & 0 & 4 \\
 7 & 9 & 1 & 0 \\
 9 & 14 & 6 & 15
 \end{bmatrix}
 =
 \begin{array}{c}
 K' \\
 \begin{array}{c}
 K'_{3,1} \\
 K'_{1,3}
 \end{array}
 \end{array}
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 884 \\
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 \end{bmatrix}$$

$$\begin{aligned}
 [20 \ 68 \ 48 \ 88] &= W_{1,-} = A_{1,-} + \phi_1 = [21 \ 67 \ 48 \ 88] + [-1 \ 1 \ 0 \ 0] \\
 [21 \ 18 \ 29 \ 42] &= W_{3,-} = A_{3,-} + \phi_3 = [22 \ 19 \ 29 \ 41] + [-1 \ -1 \ 0 \ 1]
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ϕ_1 and ϕ_3 are the random noise for $W_{1,-}$ and $W_{3,-}$, respectively

$$K'_{1,3} \neq K'_{3,1} \quad \text{but} \quad \begin{aligned}
 X_{1,3} &= (11011)_2 = f_{10,5}(876) \\
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 \end{aligned}$$

Perturbed result should be within the range

$$[(1101100000)_2, (1101111111)_2]$$

Constrained Random Perturbation

Perturbed result should be within the range

$[(1101100000)_2, (1101111111)_2]$

$$(A_{su,-}^{(t)} + \phi_{su}^{(t)}) \cdot G_{-,sv}^{(t)} \geq c_{\min}(A_{su,-}^{(t)} \cdot G_{-,sv}^{(t)}, r) \pmod{q}$$

$$(A_{su,-}^{(t)} + \phi_{su}^{(t)}) \cdot G_{-,sv}^{(t)} \leq c_{\max}(A_{su,-}^{(t)} \cdot G_{-,sv}^{(t)}, r) \pmod{q}$$

$$\phi_{su}^{(t)}(k) \in \mathbb{Z},$$

Added random perturbation

Constrained Random Perturbation

- Find constrained random perturbations
- LP Relaxation

$$(A_{s_u,-}^{(t)} + \phi_{s_u}^{(t)}) \cdot G_{-,s_v}^{(t)} \geq c_{\min}(A_{s_u,-}^{(t)} \cdot G_{-,s_v}^{(t)}, r) \pmod{q}$$

$$(A_{s_u,-}^{(t)} + \phi_{s_u}^{(t)}) \cdot G_{-,s_v}^{(t)} \leq c_{\max}(A_{s_u,-}^{(t)} \cdot G_{-,s_v}^{(t)}, r) \pmod{q}$$

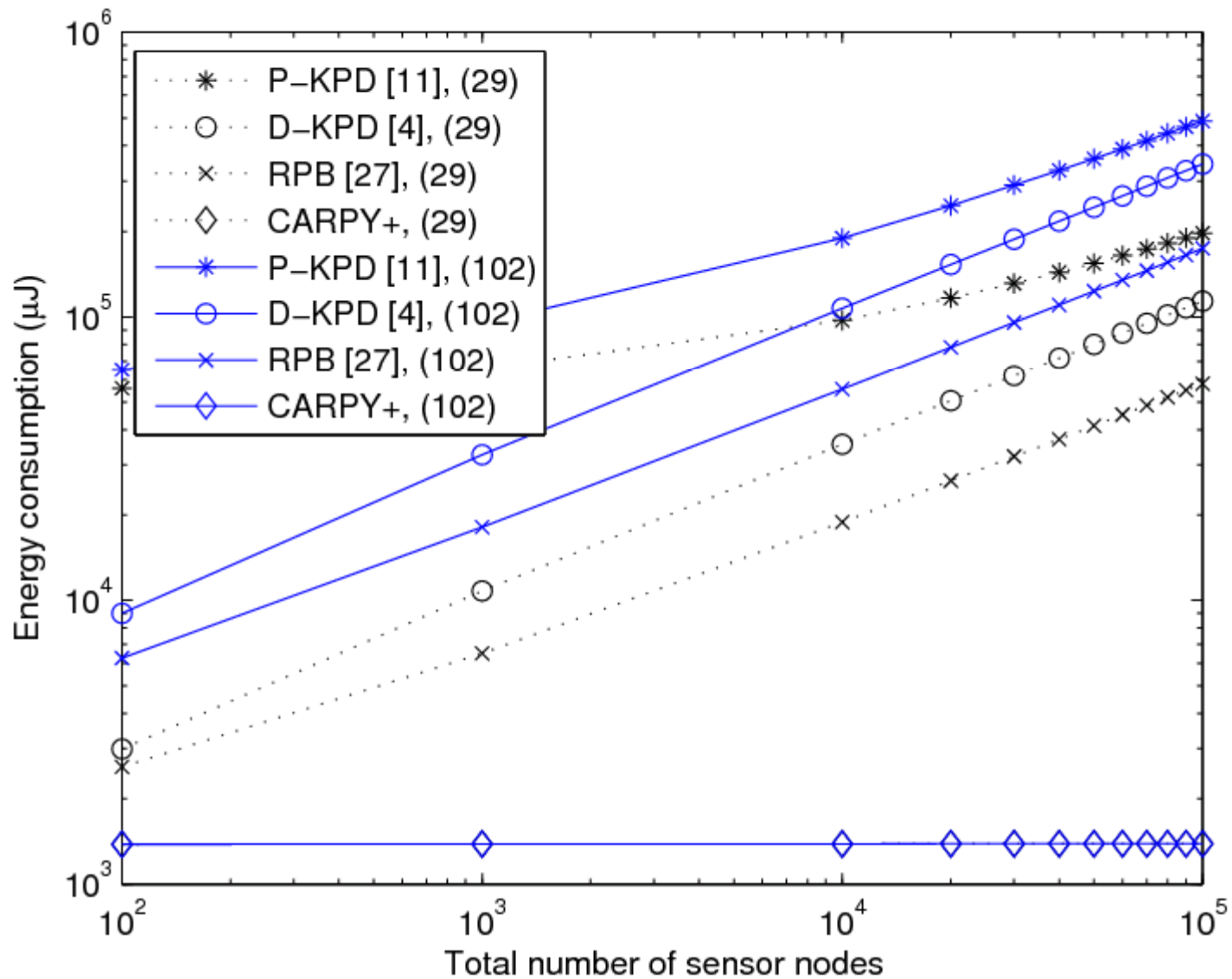
$$\phi_{s_u}^{(t)}(k) \in \mathbb{Z},$$

- Data type relaxation
 - Elements of G are float-pointed
 - randomly choose a perturbation vector, and test it

Advantage

- *All the advantages come from “non-interactive”*
 - Directed and Guaranteed Key Establishment
 - Resilience to Network Configurations
 - Efficiency
 - Resilience to Dynamic Node Deployment
 - Transparency
 - Resilience to the Adversary’s Intervention
 - *Security to node compromise*
 - *Intuitively, compromise one more node, introduce one more uncertain variable*

Energy Consumption



Prototype

- IAR Embedded Workbench, instead of TinyOS
 - Native C, instead of nesC
- TelosB (MSP430F1611, CC2420)
- 690 bytes for CODE
- 8906 bytes for DATA
- ~0.15s for key generation
- ***Our program not optimized***
 - *still working on it*

Conclusion

- A non-interactive key establishment scheme is proposed
- future work is reduce the storage cost

Thank you for your attention

For further information:

Chia-Mu Yu `r91045@csie.ntu.edu.tw`